

# Modeling switch reference in Koasati

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## 1 Introduction

- Switch reference (SR), a morphological phenomenon found in several languages in the world, is traditionally characterized as a way of indicating whether the subjects of two conjoined clauses are the same or different (Jacobsen 1993)
- Examples of SR in Koasati, a Muskogean language spoken in Louisiana and Texas, can be seen in (1)<sup>1</sup>

- (1) Joekak roomkã itcokhalihkok  
Joe-k room<sup>˜</sup> itcokhali:ka-k  
Joe-SBJ room-OBJ enter-SS  
'Joe came into the room,' (Rising 1992: 4)
- a. Edkã hihcok cokko:lit                      b. Edkã hihcan cokko:lit  
Ed<sup>˜</sup> hi:ca-k cokko:lit                      Ed<sup>˜</sup> hi:ca-n cokko:lit  
Ed-OBJ see-SS sat\_down                      Ed-OBJ see-DS sat\_down  
'saw Ed, and sat down.'                      'saw Ed, and he [Ed] sat down.'

- In (1), the morpheme *-k* (SS) in the first verb *itcokhalihkok* ('enter') indicates that its subject, *Joe*, is the same as the subject of the following verb, *hihcok/hihcan* ('see').
- In (1a), the *-k* (SS) on the second verb *hihcok* ('see') indicates that the subject of that verb, *Joe*, is the subject of the final verb *cokko:lit* ('sat down').
- In (1b), the *-n* (DS) on the second verb *hihcan* ('see') indicates that its subject, *Joe*, is not the subject of the final verb *cokko:lit* ('sat down'), but instead the object of *hihcan*, *Ed*, is.
- Consider the English equivalent of (1) in (2)

(2) Joe<sup>j</sup> came into the room. He<sub>j</sub> saw Ed<sup>k</sup>. He<sub>j/k</sub> sat down.

- *He* in the third sentence could refer to either Joe or Ed

<sup>1</sup>All data examples are copied unchanged from their sources except in the nasalization marker in examples from Kimball, which I changed from  $\check{V}$  to  $\tilde{V}$  and in the third line of the gloss. The third line of the gloss has been changed to better fit the Leipzig glossing conventions.

Gloss abbreviations: SS = SAME SUBJECT; DS = DIFFERENT SUBJECT; SBJ = SUBJECT; OBJ = OBJECT

- The English is ambiguous where the Koasati is not
- Previous semantic analyses of SR include work by Stirling (1993) and McKenzie (2007, 2011, In review)
- They analyze SR as tracking events or situations, but I pursue a reference tracking analysis for Koasati SR
- I model this data on switch reference using Predicate Logic with Anaphora (PLA; Dekker 1994), a system that maintains an ordered list of individuals in a discourse

### Roadmap

- §2 Koasati switch reference
- §3 Introduction to PLA
- §4 Initial PLA analysis: one-list system
- §5 A problem & the two-list system
- §6 Conclusion
- §A Two list fragment

## 2 Koasati switch reference

- Koasati word order is typically SOV
- SR marking appears on the verb at the end of the clause
- The verbal SS and DS morphemes are homophonous with the nominal SBJ and OBJ markings

Morpheme	Attached to Noun	Attached to Verb
<i>-k</i>	subject (SBJ)	same subject (SS)
<i>-n</i>	object (OBJ)	different subject (DS)

**Table 1:** Subject, object, and switch reference morphemes

- The overlap in the form of the nominal subject and object marker with the SR markers suggests that there is an important connection between reference and SR

### Notation for tables:

- Bold items in the table indicate overt arguments

- (1a) Joekak roomkã itcokhalihkok Edkã hihcok cokko:lit  
Joe-k room<sup>˜</sup> itcokhali:ka-k Ed<sup>˜</sup> hi:ca-k cokko:lit  
Joe-SBJ room-OBJ enter-SS Ed-OBJ see-SS sat\_down  
'Joe came into the room, saw Ed, and sat down.' (Rising 1992: 4)

Clause	Verb Gloss	Subject	Object	SR Marker
1.	entered	<b>Joe</b>	<b>room</b>	SS
2.	see	Joe	<b>Ed</b>	SS
3.	sat_down	Joe	-	-

**Table 2:** Breakdown of (1a)

- (1b) Joekak roomkã itcokhalihkok Edkã hihcan cokko:lit  
 Joe-k room-<sup>~</sup> itcokhali:ka-**k** Ed-<sup>~</sup> hi:ca-**n** cokko:lit  
 Joe-SBJ room-OBJ enter-SS Ed-OBJ see-DS sat\_down  
 ‘Joe came into the room, saw Ed, and he [Ed] sat down.’ (Rising 1992: 4)

Clause	Verb Gloss	Subject	Object	SR Marker
1.	enter	<b>Joe</b>	<b>room</b>	SS
2.	see	Joe	<b>Ed</b>	DS
3.	sat_down	Ed	-	-

**Table 3:** Breakdown of (1b)

- From these examples, it seems that there is a pattern to how individuals are introduced and referred back to
- Further, this pattern can be manipulated by the switch reference markers
  - The SS marker makes the subject and object of the SS marked clause the available subject and object, respectively, for the next clause
  - The DS marker makes the subject and object of the DS marked clause the available object and subject, respectively, for the next clause
- A system like PLA that can order individuals can be used to model this data

### 3 Background on PLA

- Predicate Logic with Anaphora (PLA; Dekker 1994) extends standard Predicate Logic in order to keep track of individuals in a discourse
- Has regular truth conditions, but a formula is interpreted as an update of an information state

(3) **A sample PLA information state**

$$s = \{ \langle a, b, c \rangle \}$$

$\uparrow$     $\uparrow$     $\uparrow$   
 $p_2$     $p_1$     $p_0$

- $p_i$ :  $i$  indexes the position of the pronoun
- $\exists$ : introduces individuals to information state

(2) Joe<sub>*j*</sub> came into the room. He<sub>*j*</sub> saw Ed<sub>*k*</sub>. He<sub>*j/k*</sub> sat down.

**Table 4:** Analysis of one translation of (2)

English	PLA	Pro. Interpr.	Output State
a.			$s_0 = \{ \langle \rangle \}$
b. Joe <sub><i>j</i></sub> came into the room.	$\exists x(x = j \wedge \exists y(y = r \wedge lxy))$		$s_1 = \{ \langle r, j \rangle \}$
c. He <sub><i>j</i></sub> saw Ed <sub><i>k</i></sub> .	$\exists y(y = e \wedge Hp_{0y})$	$[p_0]_{s_1} = j$	$s_2 = \{ \langle r, j, e \rangle \}$
d. He <sub><i>k</i></sub> sat down.	$Cp_0$	$[p_0]_{s_2} = e$	$s_3 = \{ \langle r, j, e \rangle \}$

- In (b), the narrow scope quantifier adds  $r$  to the information state first
- Then the broad scope quantifier adds  $j$  to the information state

**Table 5:** Analysis of other translation of (2)

English	PLA	Pro. Interpr.	Output State
a.			$s_0 = \{ \langle \rangle \}$
b. Joe <sub><i>j</i></sub> came into the room.	$\exists x(x = j \wedge \exists y(y = r \wedge lxy))$		$s_1 = \{ \langle r, j \rangle \}$
c. He <sub><i>j</i></sub> saw Ed <sub><i>k</i></sub> .	$\exists y(y = e \wedge Hp_{0y})$	$[p_0]_{s_1} = j$	$s_2 = \{ \langle r, j, e \rangle \}$
d. He <sub><i>j</i></sub> sat down.	$Cp_1$	$[p_1]_{s_2} = j$	$s_3 = \{ \langle r, j, e \rangle \}$

### 4 One list analysis

- In English the ambiguity of *he* is represented in PLA by different pronoun terms:  $p_0$  and  $p_1$
- The lack of ambiguity in the Koasati data can be captured by translating the subject agreement marker as  $p_0$  and object agreement marker as  $p_1$
- Further, the switch reference markers can be translated so that the DS marker swaps the order of the individuals in the  $p_0$  and  $p_1$  positions and the SS marker maintains the order
  - a-SBJ:  $\exists z(z = a)$       • intransitive verb:  $Vp_0$       • SS:  $\exists x(x = p_0 \wedge \exists y(y = p_1))$
  - b-OBJ:  $\exists x(x = p_0 \wedge \exists z(z = b))$  • transitive verb:  $Vp_0p_1$       • DS:  $\exists y(y = p_1 \wedge \exists x(x = p_0))$

(4) **SS marker**

$$s_n = \{ \langle a, b, c \rangle \} \xrightarrow{SS} s_{n+1} = \{ \langle a, b, c, b, c \rangle \}$$

(5) **DS marker**

$$s_n = \{ \langle a, b, c \rangle \} \xrightarrow{DS} s_{n+1} = \{ \langle a, b, c, c, b \rangle \}$$

- (1) Joekak roomkã itcokhalihkok  
 Joe-k room-<sup>~</sup> itcokhali:ka-**k**  
 Joe-SBJ room-OBJ enter-SS

‘Joe came into the room.’ (Rising 1992: 4)

**Table 6:** Analysis of (1)

Gloss	PLA	Pronoun Interp.	Output State
a. Joe-SBJ	$\exists z(z = j)$		$s_1 = \{ \langle j \rangle \}$
b. room-OBJ	$\exists x(x = p_0 \wedge \exists z(z = r))$	$[p_0]_{s_1} = j$	$s_2 = \{ \langle j, r, j \rangle \}$
c. enter	$!p_0p_1$	$[p_1]_{s_2} = r, [p_0]_{s_2} = j$	$s_3 = \{ \langle j, r, j \rangle \}$
d. -SS	$\exists x(x = p_0 \wedge \exists y(y = p_1))$	$[p_1]_{s_3} = r, [p_0]_{s_3} = j$	$s_4 = \{ \langle j, r, j, r, j \rangle \}$

- (1a) Edkã hihcok cokko:lit  
 Ed-<sup>~</sup> hi:ca-**k** cokko:lit  
 Ed-OBJ see-SS sat\_down  
 ‘saw Ed, and sat down.’ (Rising 1992: 4)

**Table 7:** Analysis of (1a)

Gloss	PLA	Pronoun Interp.	Output State
e. Ed-OBJ	$\exists x(x = p_0 \wedge \exists z(z = e))$	$[p_0]_{s_4} = j$	$s_5 = \{\langle j, r, j, r, j, e, j \rangle\}$
f. see	$Hp_0p_1$	$[p_1]_{s_5} = e, [p_0]_{s_5} = j$	$s_6 = \{\langle j, r, j, r, j, e, j \rangle\}$
g. -SS	$\exists x(x = p_0 \wedge \exists y(y = p_1))$	$[p_1]_{s_6} = e, [p_0]_{s_6} = j$	$s_7 = \{\langle j, r, j, r, j, e, j, e, j \rangle\}$
h. sat_down	$Cp_0$	$[p_0]_{s_7} = j$	$s_8 = \{\langle j, r, j, r, j, e, j, e, j \rangle\}$

(1b) Edkā hihcan cokko:lit  
 Ed~ hi:ca-n cokko:lit  
 Ed-OBJ see-DS sat\_down  
 ‘saw Ed, and he [Ed] sat down.’ (Rising 1992: 4)

**Table 8:** Analysis of (1b)

Gloss	PLA	Pronoun Interp.	Output State
e. Ed-OBJ	$\exists x(x = p_0 \wedge \exists z(z = e))$	$[p_0]_{s_4} = j$	$s_5 = \{\langle j, r, j, r, j, e, j \rangle\}$
f. see	$Hp_0p_1$	$[p_1]_{s_5} = e, [p_0]_{s_5} = j$	$s_6 = \{\langle j, r, j, r, j, e, j \rangle\}$
g. -DS	$\exists y(y = p_1 \wedge \exists x(x = p_0))$	$[p_1]_{s_6} = e, [p_0]_{s_6} = j$	$s_7 = \{\langle j, r, j, r, j, e, j, e \rangle\}$
h. sat_down	$Cp_0$	$[p_0]_{s_7} = e$	$s_8 = \{\langle j, r, j, r, j, e, j, e \rangle\}$

- The different SR morpheme translations in (g) for **Tables 7-8** generate distinct unambiguous interpretations

## 5 A problem

- The data in (6) cannot be accounted for using the one list system

(6) Joekak roomkā itcokhali:kon  
 Joe-k room~ itcokhali:ka-n  
 Joe-SBJ room-OBJ enter-DS  
 ‘Joe came into the room.’ (Rising 1992: 4)

a. Edkak hihcan cokko:lit  
 Ed-k hi:ca-n cokko:lit  
 Ed-SBJ see-DS sat\_down  
 ‘Ed saw him, and Joe sat down.’

Clause	Verb Gloss	Subject	Object	SR Marker
1.	enter	<b>Joe</b>	<b>room</b>	DS
2.	see	<b>Ed</b>	Joe	DS
3.	sat_down	Joe	-	-

**Table 9:** Breakdown of (6)

**Table 10:** Analysis of (6)

Gloss	PLA	Pronoun Interp.	Output State
a. Joe-SBJ	$\exists z(z = j)$		$s_1 = \{\langle j \rangle\}$
b. room-OBJ	$\exists x(x = p_0 \wedge \exists z(z = r))$	$[p_0]_{s_1} = j$	$s_2 = \{\langle j, r, j \rangle\}$
c. enter	$lp_0p_1$	$[p_1]_{s_2} = r, [p_0]_{s_2} = j$	$s_3 = \{\langle j, r, j \rangle\}$
d. -DS	$\exists y(y = p_1 \wedge \exists x(x = p_0))$	$[p_1]_{s_3} = r, [p_0]_{s_3} = j$	$s_4 = \{\langle j, r, j, j, r \rangle\}$

**Table 11:** Analysis of (6a)

Gloss	PLA	Pronoun Interp.	Output State
e. Ed-SBJ	$\exists x(x = e)$		$s_5 = \{\langle j, r, j, j, r, e \rangle\}$
f. see	$Hp_0p_1$	$[p_1]_{s_5} = r, [p_0]_{s_5} = e$	$s_6 = \{\langle j, r, j, j, r, e, r \rangle\}$
g. -DS	$\exists y(y = p_1 \wedge \exists x(x = p_0))$	$[p_1]_{s_6} = e, [p_0]_{s_6} = r$	$s_7 = \{\langle j, r, j, j, r, e, e, r \rangle\}$
h. sat_down	$Cp_0$	$[p_0]_{s_7} = r$	$s_8 = \{\langle j, r, j, j, r, e, e, r \rangle\}$

### 5.1 Two list analysis

- I adapt PLA to be a two list system
- Bittner (2001) uses a two list system for anaphora and also applies it in an analysis of the obviative system in Kalallisit (West Greenlandic) (Bittner 2011)
- Little and Moroney (2016) use a two list system related to the one presented here in an analysis of obviation in Mi'gmaq

#### (7) A sample two list information state

$$s = \{ \langle \langle a, b \rangle, \langle c, d \rangle \rangle \}$$

$$\begin{matrix} \uparrow & \uparrow & \uparrow & \uparrow \\ p_1^\top & p_0^\top & p_1^\perp & p_0^\perp \end{matrix}$$

- a-SBJ:  $\exists^\top z(z = a)$
- b-OBJ:  $\exists^\perp z(z = b)$
- intrans. verb:  $Vp_0^\top$
- trans. verb:  $Vp_0^\top p_0^\perp$
- SS:  $\exists^\perp x(x = p_0^\perp) \wedge \exists^\perp y(y = p_0^\top)$
- DS:  $\exists^\top y(y = p_0^\perp) \wedge \exists^\perp x(x = p_0^\top)$

#### (8) SS marker

$$s_n = \{ \langle \langle a, b \rangle, \langle c, d \rangle \rangle \} \xrightarrow{SS} s_{n+1} = \{ \langle \langle a, b \rangle, \langle c, d, b, d \rangle \rangle \}$$

#### (9) DS marker

$$s_n = \{ \langle \langle a, b \rangle, \langle c, d \rangle \rangle \} \xrightarrow{DS} s_{n+1} = \{ \langle \langle a, b, d \rangle, \langle c, d, b \rangle \rangle \}$$

## 5.2 Accounting for data

- The two list system can still account for the initial data:

**Table 12:** Analysis of (1)

Gloss	PLA	Pronoun Interp.	Output State
a. Joe-SBJ	$\exists^{\perp}z(z = j)$		$s_1 = \{\langle\langle j \rangle, \langle \rangle\rangle\}$
b. room-OBJ	$\exists^{\perp}z(z = r)$		$s_2 = \{\langle\langle j \rangle, \langle r \rangle\rangle\}$
c. enter	$lp_0^{\perp}p_0^{\perp}$	$[p_0^{\perp}]_{s_2} = j, [p_0^{\perp}]_{s_2} = r$	$s_3 = \{\langle\langle j \rangle, \langle r \rangle\rangle\}$
d. -SS	$\exists^{\perp}x(x = p_0^{\perp} \wedge \exists^{\perp}y(y = p_0^{\perp}))$	$[p_0^{\perp}]_{s_3} = r, [p_0^{\perp}]_{s_3} = j$	$s_4 = \{\langle\langle j \rangle, \langle r, j, r \rangle\rangle\}$

**Table 13:** Analysis of (1a)

Gloss	PLA	Pronoun Interp.	Output State
e. Ed-OBJ	$\exists^{\perp}z(z = e)$		$s_5 = \{\langle\langle j \rangle, \langle r, j, r, e \rangle\rangle\}$
f. see	$Hp_0^{\perp}p_0^{\perp}$	$[p_0^{\perp}]_{s_5} = j, [p_0^{\perp}]_{s_5} = e$	$s_6 = \{\langle\langle j \rangle, \langle r, j, r, e \rangle\rangle\}$
g. -SS	$\exists^{\perp}x(x = p_0^{\perp} \wedge \exists^{\perp}y(y = p_0^{\perp}))$	$[p_0^{\perp}]_{s_6} = j, [p_0^{\perp}]_{s_6} = e$	$s_7 = \{\langle\langle j \rangle, \langle r, j, r, e, j, e \rangle\rangle\}$
h. sat_down	$Cp_0^{\perp}$	$[p_0^{\perp}]_{s_7} = j$	$s_8 = \{\langle\langle j \rangle, \langle r, j, r, e, j, e \rangle\rangle\}$

**Table 14:** Analysis of (1b)

Gloss	PLA	Pronoun Interp.	Output State
e. Ed-OBJ	$\exists^{\perp}z(z = e)$		$s_5 = \{\langle\langle j \rangle, \langle r, j, r, e \rangle\rangle\}$
f. see	$Hp_0^{\perp}p_0^{\perp}$	$[p_0^{\perp}]_{s_5} = j, [p_0^{\perp}]_{s_5} = e$	$s_6 = \{\langle\langle j \rangle, \langle r, j, r, e \rangle\rangle\}$
g. -DS	$\exists y(y = p_0^{\perp}) \wedge \exists^{\perp}x(x = p_0^{\perp})$	$[p_0^{\perp}]_{s_6} = j, [p_0^{\perp}]_{s_6} = e$	$s_7 = \{\langle\langle j, e \rangle, \langle r, j, r, e, j \rangle\rangle\}$
h. sat_down	$Cp_0^{\perp}$	$[p_0^{\perp}]_{s_7} = e$	$s_8 = \{\langle\langle j, e \rangle, \langle r, j, r, e, j \rangle\rangle\}$

- It can account for the problematic data by keeping the available subject and object individuals separate:

**Table 15:** Analysis of (6)

Gloss	PLA	Pronoun Interp.	Output State
a. Joe-SBJ	$\exists z(z = j)$		$s_1 = \{\langle\langle j \rangle, \langle \rangle\rangle\}$
b. room-OBJ	$\exists^{\perp}z(z = r)$		$s_2 = \{\langle\langle j \rangle, \langle r \rangle\rangle\}$
c. enter	$lp_0^{\perp}p_0^{\perp}$	$[p_0^{\perp}]_{s_2} = j, [p_0^{\perp}]_{s_2} = r$	$s_3 = \{\langle\langle j \rangle, \langle r \rangle\rangle\}$
d. -DS	$\exists y(y = p_0^{\perp}) \wedge \exists^{\perp}x(x = p_0^{\perp})$	$[p_0^{\perp}]_{s_3} = r, [p_0^{\perp}]_{s_3} = j$	$s_4 = \{\langle\langle j, r \rangle, \langle r, j \rangle\rangle\}$

**Table 16:** Analysis of (6a)

Gloss	PLA	Pronoun Interp.	Output State
e. Ed-SBJ	$\exists z(z = e)$		$s_5 = \{\langle\langle j, r, e \rangle, \langle r, j \rangle\rangle\}$
f. see	$Hp_0^{\perp}p_0^{\perp}$	$[p_0^{\perp}]_{s_5} = e, [p_0^{\perp}]_{s_5} = j$	$s_6 = \{\langle\langle j, r, e \rangle, \langle r, j \rangle\rangle\}$
g. -DS	$\exists y(y = p_0^{\perp}) \wedge \exists^{\perp}x(x = p_0^{\perp})$	$[p_0^{\perp}]_{s_6} = e, [p_0^{\perp}]_{s_6} = j$	$s_7 = \{\langle\langle j, r, e, j \rangle, \langle r, j, e \rangle\rangle\}$
h. sat_down	$Cp_0^{\perp}$	$[p_0^{\perp}]_{s_7} = j$	$s_8 = \{\langle\langle j, r, e, j \rangle, \langle r, j, e \rangle\rangle\}$

## 6 Conclusion

- I have presented basic data of switch reference in Koasati
- I have discussed two PLA analyses for how to account for this data
  - One account uses Dekker's (1994) one-list system
  - The other account modifies his system to two lists to separate subjects and objects
- The two list analysis is better equipped to capture the data
- There is more work to be done to capture more complex data (plurals, indexicals, ditransitives)

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## A Two list fragment

- Additions to PLA are indicated with a \*

### DEFINITION 1.1 (Basic Expressions of PLA)

1.  $C = \{a, b, \dots, n\}$  (entity) constants
2.  $V = \{x, y, z, x', y', z', \dots\}$  (entity) variables
- \*3.  $A = \{p_i^\top \mid i \in \mathcal{N}\}$  (entity) pronouns of list  $\top$
- \*4.  $B = \{p_i^\perp \mid i \in \mathcal{N}\}$  (entity) pronouns of list  $\perp$
- \*5.  $T = C \cup V \cup A \cup B$  (entity) terms
6.  $R^n = \{A^1, \dots, A^n, B^1, \dots, Z^n\}$  n-ary predicates

### DEFINITION 1.2 (Syntax of PLA) The set $L$ of PLA formulas is the smallest set such that:

1. if  $t_1, \dots, t_n \in T$  and  $R \in R^n$ , then  $Rt_1 \dots t_n \in L$
2. if  $t_1, t_2 \in T$ , then  $t_1 = t_2 \in L$
3. if  $\phi \in L$ , then  $\neg\phi \in L$
- \*4. if  $\phi \in L$  and  $x \in V$ , then  $\exists^\top x\phi \in L$
- \*5. if  $\phi \in L$  and  $x \in V$ , then  $\exists^\perp x\phi \in L$
6. if  $\phi, \psi \in L$ , then  $(\phi \wedge \psi) \in L$

### DEFINITION 2.1 (Information States)

- \*1.  $S^n = \mathcal{P}(D^a \times D^b)$  the set of information states about  $n$  subjects, where  $a$  is the number of subject in the  $\top$  list and  $b$  is the number of subjects in the  $\perp$  list and  $a + b = n$
2.  $S = \cup_{n \in \mathcal{N}} S^n$  the set of information states
- \*3. For a state  $s \in S^n$ , where  $a + b = n$  and  $0 < j \leq a$ , and for any case  $e = \langle \langle d_1^\top, \dots, d_a^\top \rangle, \langle d_1^\perp, \dots, d_b^\perp \rangle \rangle \in s$ ,  $d_j^\top$  is a possible value for the  $j$ -th subject of  $s$ , also indicated as  $e_j^\top$ .
- \*4. For a state  $s \in S^n$ , where  $a + b = n$  and  $0 < k \leq b$ , and for any case  $e = \langle \langle d_1^\top, \dots, d_a^\top \rangle, \langle d_1^\perp, \dots, d_b^\perp \rangle \rangle \in s$ ,  $d_k^\perp$  is a possible value for the  $k$ -th subject of  $s$ , also indicated as  $e_k^\perp$ .
- \*5.  $s_0 = \{ \langle \langle \rangle, \langle \rangle \rangle \}$  (the initial state of information:  $D^0 \times D^0$ )
- \*6.  $\top^n = D^a \times D^b$  (the minimal state of information about  $n$  subjects, where  $a + b = n$ )
- \*7.  $\{e\}$  for any  $e = \langle \langle d_1^\top, \dots, d_a^\top \rangle, \langle d_1^\perp, \dots, d_b^\perp \rangle \rangle \in D^a \times D^b$  (the maximal state of information about  $n$  subjects, where  $a + b = n$ )
8.  $\perp^n = \{ \}$  (the absurd information state about  $n$  subjects, where  $n > 0$ )

### DEFINITION 2.2 (Notational Convention)

1. If  $e \in D^n$  and  $e' \in D^m$ , then  $e \cdot e' = \langle e_1, \dots, e_n, e'_1, \dots, e'_m \rangle \in D^{n+m}$
2.  $e'$  is an extension of  $e$ ,  $e \leq e'$ , iff  $\exists e'' : e' = e \cdot e''$
- \*3.  $e'$  is an extension of  $e$ ,  $e \leq e'$ , iff  $\forall e^{\top'} \in e' \exists e^{\top} \in e : e^{\top} \leq e^{\top'}$  &  $\forall e^{\perp'} \in e' \exists e^{\perp} \in e : e^{\perp} \leq e^{\perp'}$
- \*4. For  $s \in S^n$  ( $i \in D^n$ ),  $N_s = n(= a + b)$ ,  $N_a = a$ ,  $N_b = b$ , the number of subjects of  $s(i)$

### DEFINITION 2.3 (Information Update)

1. State  $s'$  is an update of state  $s$ ,  $s \leq s'$ , iff  $N_s \leq N_{s'}$ , and  $\forall e' \in s' \exists e \in s : e \leq e'$

### DEFINITION 3.1 (Interpretation of Terms)

1.  $[c]_{\mathcal{M}, s, e, g} = F(c)$  for all constants  $c$
2.  $[x]_{\mathcal{M}, s, e, g} = g(x)$  for all variables  $x$
- \*3.  $[p_i^\top]_{\mathcal{M}, s, e, g} = e_{N_\top - i}^\top$  for all pronouns  $p_i^\top$  and  $e$  and  $e^\top$  and  $s$  such that  $e^\top \in e$  and  $e \in s$  and  $N_s > i$
- \*4.  $[p_i^\perp]_{\mathcal{M}, s, e, g} = e_{N_\perp - i}^\perp$  for all pronouns  $p_i^\perp$  and  $e$  and  $e^\perp$  and  $s$  such that  $e^\perp \in e$  and  $e \in s$  and  $N_s > i$

### DEFINITION 3.2 (Semantics of PLA)

1.  $s \llbracket Rt_1 \dots t_n \rrbracket_{\mathcal{M}, g} = \{e \in s \mid \langle [t_1]_{\mathcal{M}, s, e, g}, \dots, [t_n]_{\mathcal{M}, s, e, g} \rangle \in F(R)\}$  (if  $N_s > I_{t_1, \dots, t_n}$ )
2.  $s \llbracket t_1 = t_2 \rrbracket_{\mathcal{M}, g} = \{e \in s \mid [t_1]_{\mathcal{M}, s, e, g} = [t_2]_{\mathcal{M}, s, e, g}\}$
3.  $s \llbracket \neg\phi \rrbracket_{\mathcal{M}, g} = \{e \in s \mid \neg \exists e' : e \leq e' \ \& \ e' \in s \llbracket \phi \rrbracket_{\mathcal{M}, g}\}$
- \*4.  $s \llbracket \exists^\top x\phi \rrbracket_{\mathcal{M}, g} = \{ \langle e^\top \cdot d, e^\perp \rangle \mid d \in D \ \& \ \langle e^\top, e^\perp \rangle \in s \llbracket \phi \rrbracket_{\mathcal{M}, g[x/d]} \}$
- \*5.  $s \llbracket \exists^\perp x\phi \rrbracket_{\mathcal{M}, g} = \{ \langle e^\top, e^\perp \cdot d \rangle \mid d \in D \ \& \ \langle e^\top, e^\perp \rangle \in s \llbracket \phi \rrbracket_{\mathcal{M}, g[x/d]} \}$
6.  $s \llbracket \phi \wedge \psi \rrbracket_{\mathcal{M}, g} = s \llbracket \phi \rrbracket_{\mathcal{M}, g} \llbracket \psi \rrbracket_{\mathcal{M}, g}$

### DEFINITION 4.1 (Support and Entailment)

1.  $s$  supports  $\phi$  wrt  $\mathcal{M}$  and  $g$ ,  $s \models_{\mathcal{M}, g} \phi$  iff  $\forall e \in s : \exists e' : e \leq e' \ \& \ e' \in s \llbracket \phi \rrbracket_{\mathcal{M}, g}$
2.  $\phi_1, \dots, \phi_n$  entail  $\psi$ ,  $\phi_1, \dots, \phi_n \models \psi$  iff  $\forall \mathcal{M}, g \forall s \in S : s \llbracket \phi_1 \rrbracket_{\mathcal{M}, g} \dots \llbracket \phi_n \rrbracket_{\mathcal{M}, g} \models_{\mathcal{M}, g} \psi$  (if defined)